

Syndrome Decoding with Hints

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Selected Areas in Cryptography 2026

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1 Preliminaries

- Syndrome decoding and Information Set Decoding (ISD)
- Approximate and perfect hints

2 Our contribution

- From hints to reliabilities
- Hint-ISD: biased permutations

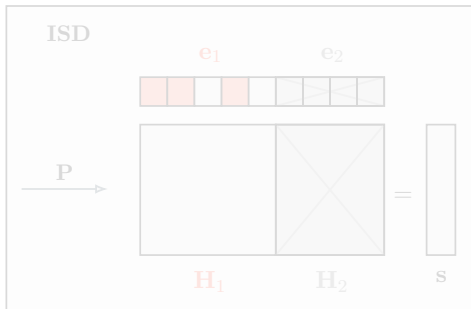
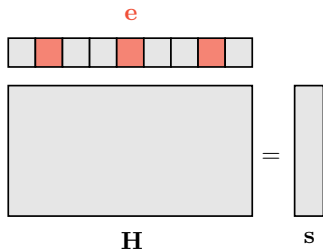
3 Experiments and practical impact on HQC and McEliece

Syndrome decoding problem and ISD

Syndrome Decoding Problem (SDP)

Given $\mathbf{H} \in \mathbb{F}_2^{(n-k) \times n}$, $\mathbf{s} \in \mathbb{F}_2^{n-k}$ and weight t , find

$$\mathbf{e} \in \mathbb{F}_2^n, \quad |\mathbf{e}| = t, \quad \mathbf{H}\mathbf{e} = \mathbf{s}.$$

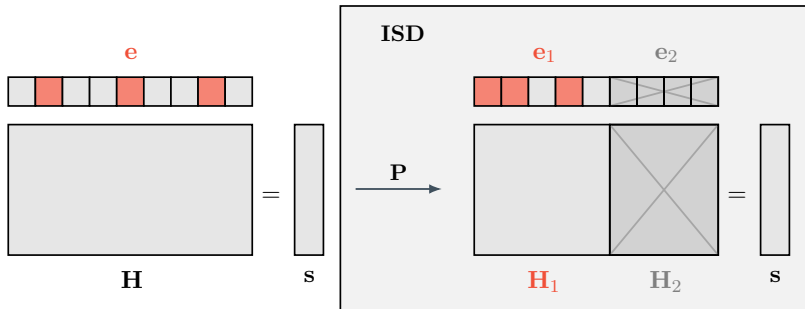


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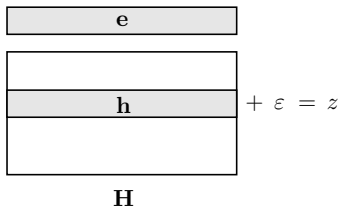


Approximate hint

For a row $\mathbf{h}$ of $\mathbf{H}$, we obtain

$$\langle \mathbf{h}, \mathbf{e} \rangle + \varepsilon = z \in \mathbb{N}, \quad \varepsilon \sim \mathcal{D}.$$

We call $(\mathbf{h}, z)$ an *approximate hint*.



Perfect hint

Special case $\varepsilon = 0$:

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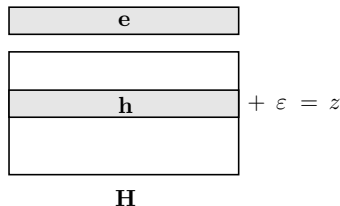


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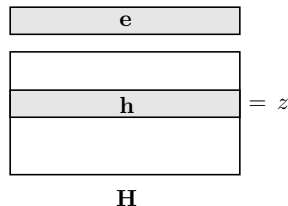


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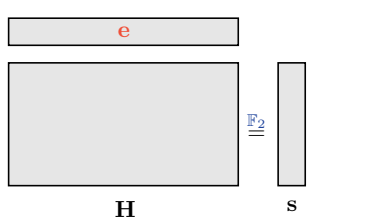
Syndrome Decoding Problem

Given

- $(\mathbf{H}, \mathbf{s})$ an instance of SDP;

Find

$$\mathbf{e} \in \mathbb{F}_2^n, \quad |\mathbf{e}| = t, \quad \mathbf{H}\mathbf{e} = \mathbf{s} \text{ on } \mathbb{F}_2,$$



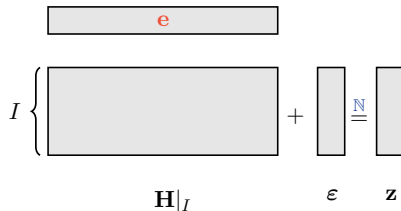
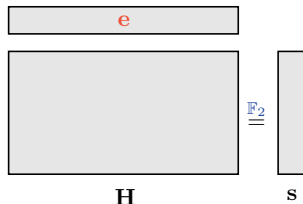
Syndrome Decoding Problem with Hints

Given

- $(\mathbf{H}, \mathbf{s})$ an instance of SDP;
- hints $\langle \mathbf{h}, \mathbf{e} \rangle + \varepsilon = \mathbf{z}$ for $\mathbf{h} \in I$, where $\varepsilon \sim \mathcal{D}$.

Find

$$\mathbf{e} \in \mathbb{F}_2^n, \quad |\mathbf{e}| = t, \quad \mathbf{H}\mathbf{e} = \mathbf{s} \text{ on } \mathbb{F}_2,$$
$$\mathbf{H}|_I \mathbf{e} + \varepsilon = \mathbf{z} \text{ on } \mathbb{N}.$$



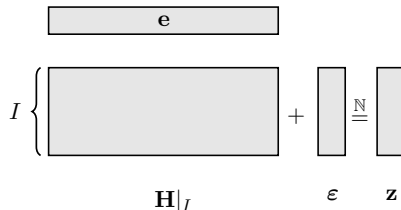
Scoring and score distribution

Goal

Use the hints

$$\mathbf{H}|_I \mathbf{e} + \boldsymbol{\varepsilon} = \mathbf{z}$$

to optimize the choice of **P**.



Score of a coordinate i [FL20]

$$\psi_i(\mathbf{z}) = \langle \mathbf{h}_i|_I, \mathbf{z} \rangle + \langle \mathbf{1}^m - \mathbf{h}_i|_I, t\mathbf{1}^m - \mathbf{z} \rangle$$

Intuition: if $i \in \text{Supp}(\mathbf{e})$, $\mathbf{h}_i|_I$ contributes to $\mathbf{z} \Rightarrow$ larger $\psi_i(\mathbf{z})$.

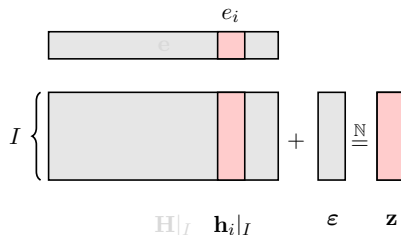
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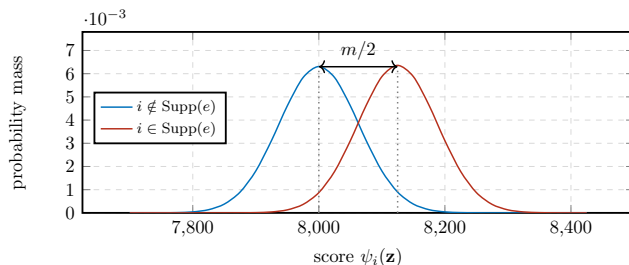
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Scoring and score distribution

Score of a coordinate i [FL20]

We actually use for $|I| = m$

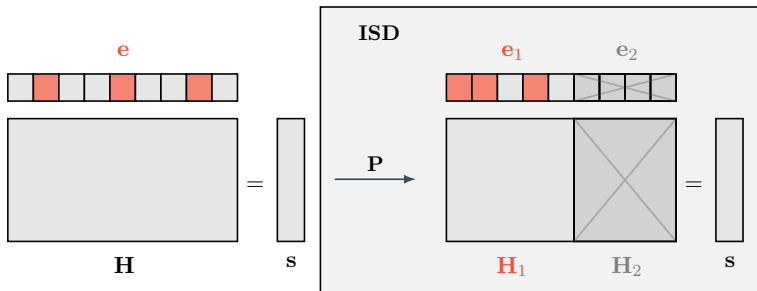
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Key idea: Larger score $\Rightarrow$ coordinate i is more likely in $\text{Supp}(\mathbf{e})$.

Integrate ISD with reliabilities: Hint-ISD

Recall: ISD chooses P at random

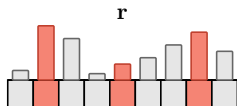


Integrate ISD with reliabilities: Hint-ISD



Reliability

$$r_i := \Pr[i \in \text{Supp}(\mathbf{e}) \mid \psi_i(\mathbf{z})]$$

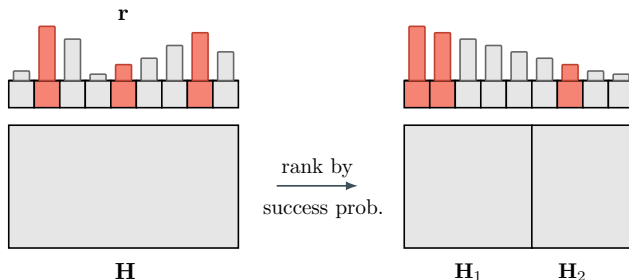


$\mathbf{H}$

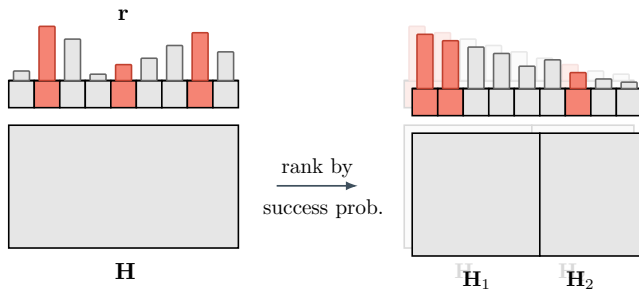


how to choose
 $\mathbf{P}$ from $\mathbf{r}$?

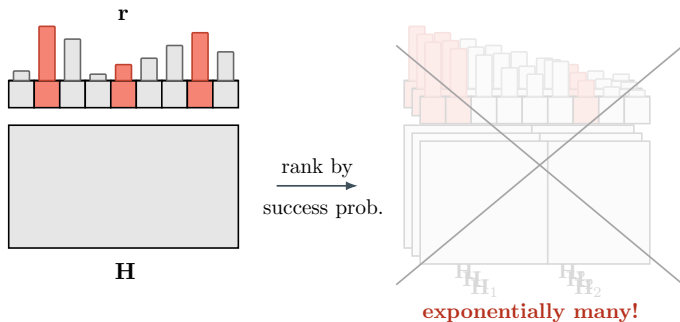
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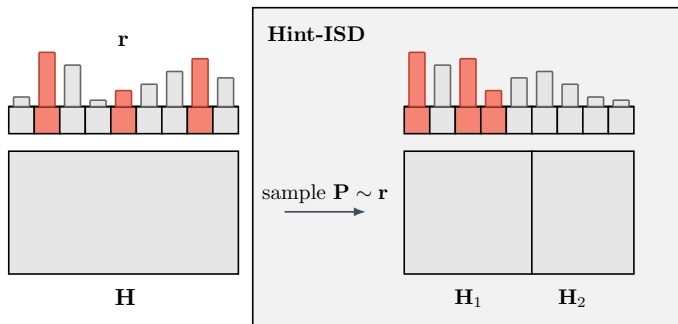


Integrate ISD with reliabilities: Hint-ISD



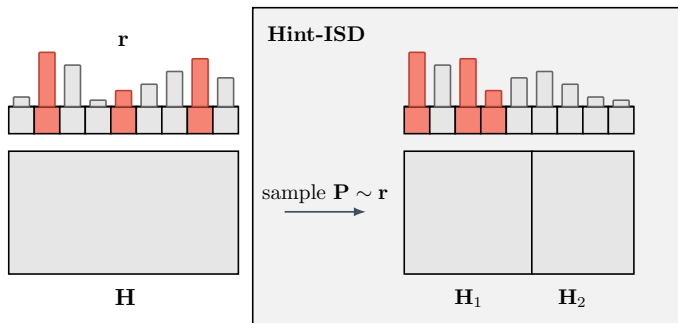
Open problem: can this be done efficiently?

Integrate ISD with reliabilities: Hint-ISD



Feasible in practice!

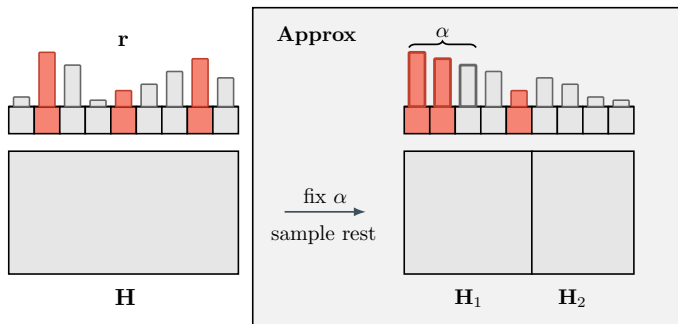
Integrate ISD with reliabilities: Hint-ISD



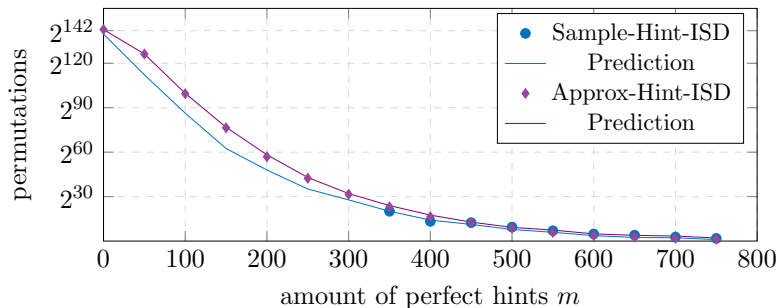
Feasible in practice!

but estimating the running time requires knowing $\mathbf{e}$

Integrate ISD with reliabilities: Hint-ISD



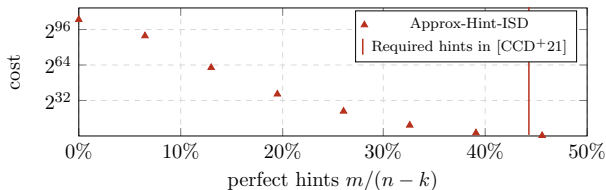
Validating approximation



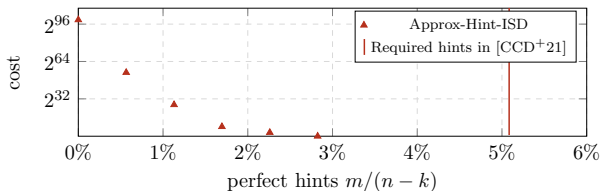
Remark: Hint-ISD can be integrated with any ISD variant.

Practical results - Perfect hints

McEliece.



HQC.



Key takeaways.

- 1 Scores can be converted into reliabilities.
- 2 Reliabilities can be integrated into ISD.
- 3 Hint-ISD gives non-trivial speedups for *any* hint exposure.
- 4 Hint-ISD interpolates between **classical ISD** and **polynomial-time decoding**.

Additional results.

- Lower bound on hints required for polynomial-time decoding.
- Stern-based Hint-ISD.
- Experiments for approximate hints.
- Analysis for systematic matrices.



[eprint.iacr.org/
2026/341](https://eprint.iacr.org/2026/341)

-  Pierre-Louis Cayrel, Brice Colombier, Vlad-Florin Drăgoi, Alexandre Menu, and Lilian Bossuet.
Message-recovery laser fault injection attack on the classic mceliece cryptosystem.
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Profiled side-channel attack on cryptosystems based on the binary syndrome decoding problem.
IEEE Transactions on Information Forensics and Security, 17:3407–3420, 2022.
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Quantitative group testing and the rank of random matrices.
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